

SURVIVAL TIME OF A RANDOM GRAPH

A. M. FRIEZE

Received September 3, 1987

Let $V_n = \{1, 2, \dots, n\}$ and e_1, e_2, \dots, e_N , $N = \binom{n}{2}$ be a random permutation of $V_n^{(2)}$. Let $E_t = \{e_1, e_2, \dots, e_t\}$ and $G_t = (V_n, E_t)$. If Π is a monotone graph property then the hitting time $\tau(\Pi)$ for Π is defined by $\tau = \tau(\Pi) = \min \{t: G_t \in \Pi\}$. Suppose now that G_t starts to deteriorate i.e. loses edges in order of age, e_1, e_2, \dots . We introduce the idea of the *survival time* $\tau' = \tau'(\Pi)$ defined by

$$\tau' = \max \{u: (V_n, \{e_u, e_{u+1}, \dots, e_N\}) \in \Pi\}.$$

We study in particular the case where Π is k -connectivity. We show that

- 1) $\lim_{n \rightarrow \infty} \Pr(\tau' \cong an) = e^{-2a}$ for $a \in \mathbb{R}^+$
- 2) $\lim_{n \rightarrow \infty} \frac{1}{n} E(\tau') = \frac{1}{n}$

i.e. τ'/n is asymptotically negative exponentially distributed with mean $\frac{1}{2}$.

1. Introduction

Let $V_n = \{1, 2, \dots, n\}$ and e_1, e_2, \dots, e_N , $N = \binom{n}{2}$ be a random permutation of $V_n^{(2)}$, the edge of set of the complete graph K_n . If $G_t = (V_n, E_t)$ where $E_t = \{e_1, e_2, \dots, e_t\}$ then the Markov chain $\mathcal{G} = (G_t)_{t=0}^N$ is the central object of study in the theory of random graphs. If Π is a monotone increasing graph property then one wishes to establish asymptotic properties of the distribution of the *hitting time*

$$\tau(\Pi) = \min \{t: G_t \in \Pi\}.$$

Erdős and Rényi [4], [5] showed this to be an interesting problem and hundreds of papers have been written on this subject since the early papers — see the recent encyclopaedic text of Bollobás [2] or the introductory text of Palmer [7]. Erdős and Rényi thought of \mathcal{G} as describing the *evolution* of some *living organism*. Suppose we pursue this analogy and think of $\tau(\Pi)$ as the time when this organism is *fully grown*. Going the way of all flesh our graph will start to *deteriorate*, say by

losing edges. We will assume that older edges disappear first. We propose to study the progress of the graph

$$H_u = (V_n, \{e_u, e_{u+1}, \dots, e_t\})$$

for various properties Π . We define the *survival time* $\tau(\Pi)$ for the process \mathcal{G} by

$$\tau'(\Pi) = \max \{u: H_u \in \Pi\}.$$

The property discussed in this paper is k -connectivity. One can imagine many other properties worthy of study and we list some open problems at the end of the paper.

Our main result can be expressed as follows: let $k \geq 1$ be a fixed integer. Let D_k denote the property of having minimum degree at least k and let C_k denote the property of being k -connected

Theorem

If Π is C_k or D_k then

- (i) $\lim_{n \rightarrow \infty} \Pr(\tau'(\Pi) \geq an) = e^{-2a}$ for $a \in \mathbb{R}^+$
- (ii) $\lim_{n \rightarrow \infty} \frac{1}{n} E(\tau'(\Pi)) = \frac{1}{2}$
- (iii) $\lim_{n \rightarrow \infty} \Pr(\tau'(C_k) = \tau'(D_k)) = 1. \quad \blacksquare$

2. Preliminaries

We need some probability inequalities related to the k -connectivity of a random graph. The main thrust of the assertions are from Erdős and Rényi [6] but the calculations giving the precise bounds used are left to an appendix.

Let $Z_m^{(r)}$ be the set of vertices of degree r in G_m and let $z_m^{(r)} = |Z_m^{(r)}|$.

Let $S \subseteq V_n$ be a *non-trivial* separator of G_m if $G_m[V_n - S]$ is not connected but has no isolated vertices.

Lemma 1

Let

$$(2.1) \quad m = \frac{1}{2} n \log n + \frac{1}{2} (k-1) n \log \log n - \frac{1}{2} \omega n \quad \text{be integer.}$$

(a) Suppose that in (2.1) $\omega = \omega(n) \rightarrow \infty$ but $\omega = o(\log \log n)$. In what follows $\alpha = \alpha(k)$ is some 'constant' (depending only on the constant k) whose exact value is unimportant.

$$(2.2a) \quad \Pr(\delta(G_m) < k-1) \leq \frac{\alpha e^\omega}{\log n}$$

where $\delta(G)$ denotes the minimum degree of graph G .

$$(2.2b) \quad \Pr(\delta(G_m) \geq k) \leq \alpha e^{-\omega}$$

$$(2.2c) \quad \Pr\left(\left|z_m^{(k-1)} - \frac{e^\omega}{(k-1)!}\right| \geq \frac{\varepsilon e^\omega}{(k-1)!}\right) \leq \frac{\alpha e^{-\omega}}{\varepsilon^2} \quad \text{for } 0 < \varepsilon < 1$$

$$(2.2d) \Pr(G_m \text{ has a non-trivial separator of size } \leq k-1) \leq \alpha \frac{(\log n)^2}{\sqrt{n}}$$

$$(2.2e) \Pr(\exists \omega: |\omega| \leq \log \log n, m \text{ (as in (2.1)) such that } G_m \text{ has a non-trivial separator } S \text{ of size } \leq k-1 \text{ for which } G_m[V_n - S] \text{ has a component of size } t, 3 \leq t \leq \frac{1}{2}|V_n - S|) = O\left(\frac{(\log n)^2}{\sqrt{n}}\right).$$

$$(2.2f) \Pr(\exists v, w \in Z_m^{(k-1)}: \text{the distance from } v \text{ to } w \text{ in } G_m \text{ is 2 or less}) \leq \alpha \sqrt{\frac{\log n}{n}}.$$

(b) If we relax the restriction $\omega = o(\log \log n)$ in (a) to $\omega n \leq m$ then we can still prove

$$(2.3) \Pr(\delta(G_m) \geq k) \leq \alpha e^{-\omega/2}.$$

(c) Let now $m = \frac{1}{2}n \log n + \frac{1}{2}(k-1)n \log \log n + \frac{1}{2}\omega n$ where $\omega = \omega(n) \rightarrow \infty$. Then

$$(2.4) \Pr(\delta(G_m) < k) \leq \max\{\alpha e^{-\omega}, n^{-2}\}. \quad \blacksquare$$

The hitting time of C_k was discussed by Bollobás and Thomason [3]. That paper established

$$(2.5) \lim_{n \rightarrow \infty} \Pr(\tau(C_k) = \tau(D_k)) = 1.$$

Let $m_1 = \left\lfloor \frac{1}{2}n \log n + \frac{1}{2}(k-1)n \log \log n - \frac{1}{2}n \log \log \log n \right\rfloor$ and $m^* = \tau(D_k)$. It will be useful to think of G_m^* in the following terms (Bollobás [1]):

$$E_{m^*} = E_{m_1} \cup X \cup Y$$

where

$$X = \{e \in E_{m^*} - E_{m_1}: e \cap Z_{m_1}^{(k-1)} = \emptyset\}$$

and

$$Y = E_{m^*} - (E_{m_1} \cup X).$$

This is valid, conditional on an event of probability $1 - O((\log \log n)^{-1})$ — see (2.2b) with $m = m_1$ and $\omega = \log \log \log n$.

Now, in this case,

$$(2.6) \quad X \text{ is a random } |X| \text{-subset of } (V_n - Z_{m_1}^{(k-1)})^{(2)}.$$

Also

$$(2.7) \quad \Pr(|X \cup Y| > n \log \log \log n) \leq \Pr(\delta(G_m) \leq k-1) \leq \frac{\alpha}{\log \log n}.$$

where $m_2 = m_1 + [n \log \log \log n]$, by (2.4). Furthermore

$$(2.8) \quad \Pr(\exists e \in Y: e \subset Z_{m_1}^{(k-1)}) \leq \frac{\log n}{n}$$

and

$$(2.9) \quad \Pr(\exists z \in Z_{m_1}^{(k-1)}: |\{e \in Y: z \in e\}| > 6 \log \log \log n) \leq \frac{1}{\log \log n}.$$

Given $|X \cup Y| \leq n \log \log \log n$ and $|Z_{m_1}^{(k-1)}| = O(\log \log n)$ (use $\varepsilon=1$ in (2.2c)) it is unlikely that the conditions in (2.8), (2.9) will be violated — we are after all adding at most $n \log \log \log n$ random edges.

3. Proof of the Theorem

(i) and (iii).

Let $a \in \mathbb{R}^+$ and $u = [an]$. We show first that

$$(3.1) \quad \lim_{n \rightarrow \infty} \Pr(\delta(H_u) = k) = e^{-2a}.$$

We aim in fact to prove

$$(3.2a) \quad |\Pr(\delta(H_u) = k) - e^{-2a}| = O\left(\frac{(\log \log \log n)^2}{\log \log n}\right)$$

where the hidden constant in the “big O ” notation may depend on k . Thus from now on, if after describing an event \mathcal{E} we write $[P=1-o(1)]$, we mean $\Pr(\mathcal{E}) = 1 - O((\log \log \log n)^2 / \log \log n)$.

For the proof of (i) and (iii) we only require that a be a constant. However to prove (ii) we need to allow a to grow with n . In what follows we will assume only that

$$(3.2b) \quad a \leq \log \log \log \log n.$$

Let now $Z^{(k-1)} = Z_{m_1}^{(k-1)}$ and $\hat{Z}^{(k-1)}$ be the set of vertices of degree $k-1$ in the graph $H' = (V_n, \{e_u, e_{u+1}, \dots, e_{m_1}\})$. Note that H' has the same distribution as G_{m_1-u+1} and we may use Lemma 1 (a) with $\omega = \log \log \log n + 2a + O(1/n)$. (The $O(1/n)$ term accounts for ω_n integral.)

G_{m_1} is obtained from H' by adding $u-1$ random edges and so $Z^{(k-1)} \subseteq \hat{Z}^{(k-1)}$. By applying (2.2c) twice: once with $\omega = \log \log \log n$ and $\varepsilon = \omega^{-1}$ and once with $\omega = \log \log \log n + 2a$, ε as before, we obtain

$$(3.3) \quad \Pr\left(\frac{1-\varepsilon}{1+\varepsilon} \leq e^{2a} \frac{|Z^{(k-1)}|}{|\hat{Z}^{(k-1)}|} \leq \frac{1+\varepsilon}{1-\varepsilon}\right) \geq 1 - \frac{2a(\log \log \log n)^2}{\log \log n}.$$

Now let $G'_m = (V_n, E'_m)$ where $E'_m = E_m - E_{u-1}$ and note that G'_m has the same distribution as G_{m-u+1} . Let $\hat{m} = \min\{m: \delta(G'_m) \geq k\}$, and let \hat{v} be the unique $[P=1-o(1)]$, see (2.8)) vertex of degree $k-1$ in $G'_{\hat{m}-1}$. Then

$$\delta(H_u) = k \leftrightarrow \hat{v} \in Z^{(k-1)}.$$

Now \hat{v} is “close to” being a random element of $\hat{Z}^{(k-1)}$ and so we can see from (3.3) that $\Pr(\hat{v} \in Z^{(k-1)}) \approx e^{-2a}$, but let us do this more carefully.

Consider now the set $\{e_{m_1}, e_{m_1+1}, \dots, e_N\}$ and in particular the subset F of edges incident with one vertex in $\hat{Z}^{(k-1)}$ and one vertex not in $\hat{Z}^{(k-1)}$. We know $[P=1-o(1)]$, see (2.8) that $e_n \in F$. For $v \in \hat{Z}^{(k-1)}$ let d_v be the degree of v in G_{m_1} . We work on the assumption that $\hat{Z}^{(k-1)}$ is an independent set in G_{m_1} , $[P=1-o(1)]$. Now $d_v = k-1$ for $v \in Z^{(k-1)}$ and $d_v \leq 6 \log \log n$ otherwise $[P=1-o(1)]$, see (2.9). If d_v were the same for all $v \in \hat{Z}^{(k-1)}$ then we could deduce that $\Pr(\hat{v} \in Z^{(k-1)}) = |Z^{(k-1)}|/|\hat{Z}^{(k-1)}|$ and we would be done. This is nearly so and for each $v \in \hat{Z}^{(k-1)}$

we randomly select a set $F_v \subseteq F$ of $n_1 = [n - 1 - 6 \log \log n]$ edges incident with v . Let $F' = \bigcup_{v \in Z^{(k-1)}} F_v$. Then

$$\Pr(\hat{v} \in Z^{(k-1)} | e_{\hat{m}} \in F') = \frac{|Z^{(k-1)}|}{|\hat{Z}^{(k-1)}|}$$

and

$$\Pr(e_{\hat{m}} \notin F') = \frac{1}{n_1} + O\left(\frac{\log n}{n}\right).$$

(The $O\left(\frac{\log n}{n}\right)$ term above, accounts for $e_{\hat{m}} \in \hat{Z}^{(k-1)}$). This completes the proof of (3.2a).

We show next that

$$(3.4) \quad \lim_{n \rightarrow \infty} \Pr(\tau'(D_k) = \tau'(C_k)) = 1.$$

Observe that if $\tau'(C_k) < u = \tau'(D_k)$ then either $\tau(D_k) \neq \tau(C_k)$ or there exists $S \subseteq V_n$, $|S| = k-1$ such that

$$(3.5) \quad S \text{ is a non-trivial separator of } H_{u-1}.$$

Note next that (3.2) implies

$$\lim_{n \rightarrow \infty} \Pr(\tau'(D_k) \cong n \log \log \log \log n) = 0.$$

Thus if $u_1 = [n \log \log \log \log n]$, $\tau'(C_k) < u < u_1$, $K = (V_n, e_{u_1}, e_{u_1+1}, \dots, e_{m_1})$ and S is as in (3.5), then either

(a) S is a non-trivial separator of H_{u-1} and $H_{u-1}[V_n - S]$ has a component of size t , $3 \leq t \leq \frac{1}{2} |V_n - S|$,

or

(b) S is a non-trivial separator of K

or

(c) $\delta(K) \leq k-2$,

or

(d) 2 vertices of K of degree $k-1$ share a common neighbour.

But K has the same distribution as $G_{m_1 - u_1 + 1}$ and (2.2) shows that these 4 events all have probability tending to zero.

(i) and (iii) follow directly from (3.1) and (3.4).

(ii).

We use

$$E\left(\frac{1}{n} \tau'(\Pi)\right) = \int_0^{(1/2)n} \Pr(\tau'(\Pi) \geq nx) dx$$

and

$$\tau'(D_k) \cong \tau'(C_k) \text{ whenever } \tau(D_k) = \tau(C_k).$$

Let $\lambda = \log \log \log \log n$ and \mathcal{E} denote the event " $\tau(D_k) = \tau(C_k)$ ". Then

$$(3.6) \quad \Pr(\bar{\mathcal{E}}) = O\left(\frac{(\log n)^2}{\sqrt{n}}\right).$$

This follows from (2.2d) and the proof of Theorem VII. 4 of [2]. Now

$$(3.7) \quad \begin{aligned} E\left(\frac{1}{n} \tau'(C_k)\right) &\cong \int_0^\lambda \Pr(\tau'(C_k) \cong nx) dx \cong \\ &\cong \int_0^\lambda e^{-2x} dx - O\left(\frac{\lambda(\log \log \log n)^2}{\log \log n}\right) = \text{by (3.2)} \\ &= \frac{1}{2} - o(1). \end{aligned}$$

Now a lower bound for $E(\tau'(D_k))$.

$$\begin{aligned} E\left(\frac{1}{n} \tau'(D_k)\right) &= E\left(\frac{1}{n} \tau'(D_k) | \mathcal{E}\right) \Pr(\mathcal{E}) \cong \\ &\cong E\left(\frac{1}{n} \tau'(C_k) | \mathcal{E}\right) \Pr(\mathcal{E}) = \\ &= E\left(\frac{1}{n} \tau'(C_k)\right) - E\left(\frac{1}{n} \tau'(C_k) | \bar{\mathcal{E}}\right) \Pr(\bar{\mathcal{E}}) \cong \\ &\cong \frac{1}{2} - o(1) - E\left(\frac{1}{n} \tau(C_k) | \bar{\mathcal{E}}\right) \Pr(\bar{\mathcal{E}}). \end{aligned}$$

But

$$E\left(\frac{1}{n} \tau(C_k) | \bar{\mathcal{E}}\right) \cong (\log n)^2 + n \Pr(\tau(C_k) \cong n(\log n)^2) / \Pr(\bar{\mathcal{E}}).$$

Thus, using (3.6) and $\Pr(\tau(C_k) \cong n(\log n)^2) = o(1/n)$ (very crudely), we have

$$(3.8) \quad E\left(\frac{1}{n} \tau(C_k) | \bar{\mathcal{E}}\right) \Pr(\bar{\mathcal{E}}) = o(1)$$

and so

$$E\left(\frac{1}{n} \tau'(D_k)\right) \cong \frac{1}{2} - o(1).$$

Now for upper bounds:

$$(3.9) \quad \begin{aligned} E\left(\frac{1}{n} \tau'(D_k)\right) &= \int_0^\lambda \Pr(\tau'(D_k) \cong nx) dx + \int_\lambda^{(1/2)n} \Pr(\tau'(D_k) \cong nx) dx \cong \\ &\cong \int_0^\lambda e^{-2x} dx + O\left(\frac{\lambda(\log \log \log n)^2}{\log \log n}\right) + \int_\lambda^\infty 2\alpha e^{-x/2} dx + \int_\lambda^{(1/2)n} n^{-2} dx, \\ &= \frac{1}{2} + o(1). \end{aligned}$$

The first integral in (3.9) is approximated as in (3.7). For the second note that

$$(3.10) \quad \Pr(\tau'(D_k) \cong nx) \leq \Pr(\delta(G_{m^+}) < k) + \Pr(\delta(G_{m^-}) \cong k)$$

where

$$m^+ = \frac{1}{2} n \log n + \frac{1}{2} (k-1) n \log \log n + \frac{1}{2} nx$$

and

$$m^- = m^+ - nx.$$

Now use (2.3) with $m=m^-$ and (2.4) with $m=m^+$ in (3.10). Now an upper bound for $E(\tau'(c_k))$.

$$\begin{aligned} E\left(\frac{1}{n} \tau'(C_k)\right) &= E\left(\frac{1}{n} \tau'(C_k) | \mathcal{E}\right) \Pr(\mathcal{E}) + E\left(\frac{1}{n} \tau'(C_k) | \bar{\mathcal{E}}\right) \Pr(\bar{\mathcal{E}}) \leq \\ &\leq E\left(\frac{1}{n} \tau'(D_k) | \mathcal{E}\right) \Pr(\mathcal{E}) + E\left(\frac{1}{n} \tau(C_k) | \bar{\mathcal{E}}\right) \Pr(\bar{\mathcal{E}}) \leq \\ &\leq E\left(\frac{1}{n} \tau'(D_k)\right) + o(1) \leq \quad (\text{by (3.8)}) \\ &\leq \frac{1}{2} + o(1) \quad (\text{by (3.10)}). \end{aligned}$$

This completes the proof of our theorem.

4. Revival Time

We gain some insight into the distribution of survival time by considering the *revival time*. Let $u=\tau'(H)$ and consider adding random edges to H_u until a graph with property Π is obtained. The number of edges added $\tau''(\Pi)$ is called the revival time. It is straightforward to show that we can replace τ' by τ'' in our theorem and obtain a valid result.

For example: if $\Pi=D_k$ then $H_u[P=1-o(1)]$ contains a unique vertex v of minimal degree $k-1$. $\tau'' \cong an$ if v is not incident with any of the first an random edges added to H_u . The probability of this is approximately $\left(1 - \frac{2}{n}\right)^{an} \approx e^{-2a}$. For $\Pi=C_k$ we use the smaller likelihood of non-trivial separators.

It seems now that in general one may be able to guess the result for survival time by computing the revival time, which seems easier. One then needs a few asymptotic calculations for verification.

5. Open Questions

What is the survival time for the following properties:

- (1) having a cycle?
- (2) having a cycle of size k ?
- (3) having a path of length k ?
- (4) having a tree component of size k ?
- (5) having a clique of size k ?
- (6) having any fixed subgraph?
- (7) being non-planar?
- (8) having diameter k ?
- (9) having a vertex of degree k ?

One can also consider similar problems for random bipartite graphs, digraphs or subgraphs of the n -cube.

There are 2 conspicuous omissions from our list. These are perfect matchings and hamilton cycles. If H_k is the property of having $\lfloor k/2 \rfloor$ edge disjoint hamilton cycles plus a further edge disjoint matching of size $\lfloor n/2 \rfloor$, if k is odd, then it seems fairly clear that we can add H_k to our theorem. The proof does not require any new ideas and would be rather long, too long for this paper.

Appendix

As usual let G_p , $p = \frac{N}{m}$ denote the random graph in which edges are independently included with probability p . Let Π be any graph property. Then

$$(A0) \quad \Pr(G_p \in \Pi) = \sum_{m'} \Pr(G_{m'} \in \Pi) \Pr(G_p \text{ has } m' \text{ edges}).$$

Thus if Π is monotone i.e. if it is preserved either by adding edges or by deleting edges, then, for large n

$$(A1) \quad \Pr(G_m \in \Pi) \leq 3 \Pr(G_p \in \Pi).$$

We can thus work mainly with G_p and multiply our estimates by 3. Our inequalities are only required to hold for n large.

Proof of (2.2a).

For $k \geq 2$

$$\begin{aligned} \Pr(\delta(G_p) \leq k-2) &\leq n \sum_{t=0}^{k-2} \binom{n-1}{t} p^t (1-p)^{n-1-t} \leq \\ &\leq 2n \sum_{t=0}^{k-2} \frac{n^t}{t!} \left(\frac{\log n}{n} \right)^t \frac{(\log n)^{-(k-1)} e^\omega}{n} \leq \\ &\leq \frac{3e^\omega}{(k-2)! \log n} \end{aligned}$$

Now use (A1).

Proof of (2.2b).

Let $z_p^{(k-1)}$ be the number of vertices of degree $k-1$ in G_p .

$$\begin{aligned} E(z_p^{(k-1)}) &= n \binom{n-1}{k-1} p^{k-1} (1-p)^{n-k} = \\ &= \left(1 + O\left(\frac{(\log n)^2}{n}\right) \right) \left(\frac{np}{\log n} \right)^{k-1} \frac{e^\omega}{(k-1)!} = \\ &= \left(1 + O\left(\frac{\log \log n}{\log n}\right) \right) \frac{e^\omega}{(k-1)!}. \end{aligned}$$

Preparing for the Chebyshev inequality

$$\begin{aligned} &E(z_p^{(k-1)}(z_p^{(k-1)} - 1)) = \\ &= n(n-1) \left(p \left(\binom{n-2}{k-2} p^{k-2} (1-p)^{n-k} \right)^2 + (1-p) \left(\binom{n-2}{k-1} p^{k-1} (1-p)^{n-k-1} \right)^2 \right) = \\ &= n(n-1) \left(p \left(\frac{E(z_p^{(k-1)})}{n(n-1)p} \right)^2 + (1-p) \left(\frac{E(z_p^{(k-1)})}{n(1-p)} \frac{n-1}{n-k} \right)^2 \right) = \\ &\leq (1+2p)(E(z_p^{(k-1)}))^2 \end{aligned}$$

and so

$$\begin{aligned} \text{(A2)} \quad \text{Var}(z_p^{(k-1)}) &\leq E(z_p^{(k-1)}) + 2pE(z_p^{(k-1)})^2 \leq \\ &\leq 2E(z_p^{(k-1)}) \quad \text{as } \omega = o(\log \log n). \end{aligned}$$

Thus by the Chebyshev inequality

$$\Pr(z_p^{(k-1)} = 0) \leq \frac{2}{E(z_p^{(k-1)})} \leq 3(k-1)! e^{-\omega}.$$

Thus

$$\Pr(\delta(G_p) \geq k) \leq 3(k-1)! e^{-\omega}.$$

We can now use (A1).

Proof of (2.2c).

The property in question is not monotone (or even convex) [2] and so we cannot use (A1). We can however assert using (A2) that

$$\Pr\left(\left|z_p^{(k-1)} - \frac{e^\omega}{(k-1)!}\right| \geq \frac{\varepsilon e^\omega}{(k-1)!}\right) \leq \frac{2e^{-\omega}}{\varepsilon^2}.$$

It now follows from (A0) that there exists $m', m - \sqrt{n \log n} \leq m' \leq n$, such that

$$\Pr\left(\left|z_{m'}^{(k-1)} - \frac{e^\omega}{(k-1)!}\right| \geq \frac{\varepsilon e^\omega}{(k-1)!}\right) \leq \frac{2e^{-\omega}}{\varepsilon^2}.$$

As in the proof of (2.2a) we can deduce that

$$\Pr(\delta(G_{m'}) < k-1) \leq \frac{10e^\omega}{(k-2)!\log n}.$$

Since G_m is obtained from $G_{m-m'}$ by adding $m-m'$ random edges, we have that given $\mathcal{E} = \{\delta(G_{m'}) \leq k-1 \text{ and } z_{m'}^{(k-1)} \text{ is 'close' to its mean}\}$,

$$\begin{aligned} \Pr(z_m^{(k-1)} \neq z_{m'}^{(k-1)} | \mathcal{E}) &= O\left(\frac{\sqrt{n \log n n e^\omega}}{N-n}\right) = \\ &= O\left(\sqrt{\frac{\log n}{n}} e^\omega\right) \end{aligned}$$

and (2.2c) follows.

Proof of (2.2d).

Let θ be the probability that G_p has a non-trivial separator of size s , $0 \leq s \leq k-1$. Then

$$\theta \leq \sum_{s=0}^{k-1} \sum_{t=2}^{(1/2)(n-s)} \binom{n}{s} \binom{n-s}{t} t^{t-2} p^{t-1} (1-(1-p)^t)^s (1-p)^{t(n-s-t)}$$

(choose an s -set S for the separator, a t -set T for a small component, a spanning tree of T . Multiply by the probability that the edges of the tree exist and there is at least one v, T edge for each $v \in S$ and no $S, S \cup T$ edges.)

Thus

$$\begin{aligned} \theta &\leq \sum_{s=0}^{k-1} \sum_{t=2}^{(1/2)(n-s)} \left(\frac{ne}{s}\right)^s \left(\frac{ne}{t}\right)^t t^{t-2} p^{t-1} (tp)^s e^{-t(n-s-t)p} \leq \\ (A3) \quad &\leq 8 \sum_{s=0}^{k-1} \sum_{t=2}^{(1/2)(n-s)} \frac{1}{p} (npe^{-np} t^{s-2/t} e^{(s+t+1)p})^t = \\ &= O\left(\frac{\log n}{n}\right) \end{aligned}$$

(for small t , the "complex" term above is $O\left(\left(\frac{e^\omega \log n}{n}\right)^t\right)$. For larger t , it is $O\left(\left(\frac{e^\omega \log n}{\sqrt{n}}\right)^t\right)$. Unfortunately, we are not dealing with a monotone property. However (A0) implies

$$\Pr(G_m \text{ has a non-trivial separator}) \leq \theta \Pr(G_p \text{ has } m \text{ edges})^{-1}$$

and (2.2d) follows. (See [2] Theorem II.2.)

Proof of (2.2e).

We only have to consider the sum in (A3) for $t \geq 3$, which is then $O\left(\left(\frac{\log n}{n}\right)^2\right)$. We only have to multiply this by $O\left(n \log \log n \times \frac{\sqrt{n}}{\sqrt{\log n}}\right)$ in order to account for the number of different values of m and the transition from G_p to G_m .

Proof of (2.2f).

We can clearly assume $k \geq 2$. Let ψ be the probability that there exist $v, w \in Z_m^{(k-1)}$ at distance 2 or less from each other in G_p . Then

$$\psi \leq \sum_{r=2}^3 \frac{r!}{2} \binom{n}{r} p^{r-1} \left(\binom{n-r}{k-2} p^{k-2} (1-p)^{n-k-r+2} \right)^2 + \\ + 3 \binom{n}{3} p^3 \left(\binom{n-3}{k-2} p^{k-2} (1-p)^{n-k-1} \right)^2$$

(the first term above deals with paths of length $r=2$ or 3 and the second term deals with triangles.) Thus $\psi = O\left(\frac{1}{n}\right)$ and we can finish as in the proof of (2.2d).

Proof of (2.3).

The proof used for (2.2b) will be valid for $p \geq p_0 = \frac{\log n}{4n}$, say. All that is needed for smaller p is to show

$$\Pr(G_{p_0} \text{ has no isolated vertices}) = O\left(\frac{e^{np}}{n}\right).$$

This can be done using the Chebycheff inequality.

Proof of (2.4).

Non-trivial separators are handled as in the proof of (2.2d). The minimum degree calculation only requires the use of the expected number of vertices of degree $k-1$ or less. ■

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A. M. Frieze

Department of Mathematics
Carnegie Mellon University
Pittsburgh, PA 15213